

RTC-Tools 2.0: Waar staan we nu?

Jorn Baayen

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Outline

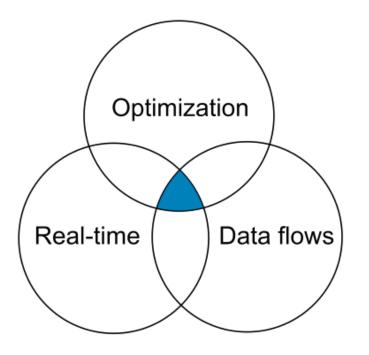
- RTC-Tools: Scope, history
- Convex optimization
- Hydraulic modelling
- Multiple objectives
- Mixed integer optimization

RTC-Tools: Scope

RTC-Tools is the Deltares toolbox for control and optimization of environmental systems.

Delft-FEWS is an open data handling platform, used for the aggregation of (real-time) environmental data flows.

Together, they provide a platform for the development of decision support systems.



History

- 2005: Reservoir module for Delft-FEWS.
- 2012: Dirk Schwanenberg releases first version of RTC-Tools source code to the public. RTC-Tools 1.x connected non-linear hydraulic and reservoir models to the IPOPT optimizer.
 - Promising results, many scientific publications
 - High interest from reservoir operators
 - But challenging to operationalize, and hard to extend
- 2015: Work starts on new mathematically rigorous foundation, initially as an experiment of Jorn Baayen and Matthijs den Toom.

- 2016: First pilot project on new foundation. Peter Gijsbers develops water allocation tool for Rijkswaterstaat using new framework.
- 2016: RTC-Tools 2.0 released.

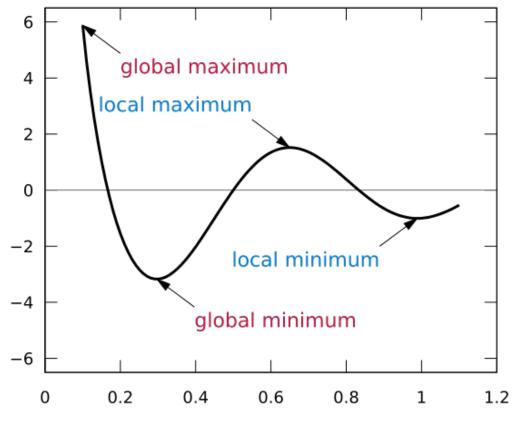
A decision support system that is used day in, day out needs to be reliable. This need can be made precise with six axioms:

- Robustness: The solutions are robust in the sense that system constraints are satisfied regardless of the differences between model and reality.
- *Feasibility*: A feasible solution always exists.
- *Quality*: Any solution is a "good" solution.
- *Stability*: The solutions are stable in the sense that small perturbations in the configuration result in small changes in the solution.
- *Determinism*: Given the same initial solution guess and configuration, the solution is always identical.

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• *Bounded solution time*: A solution is found within a predetermined amount of time.

Local and global optima



Source: Wikipedia. GFDL 1.2.

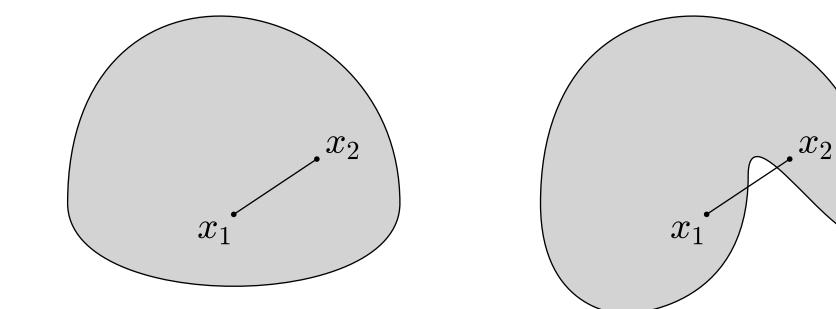
Suppose we had an optimization problem that would only have globally optimal solutions.

That would give us:

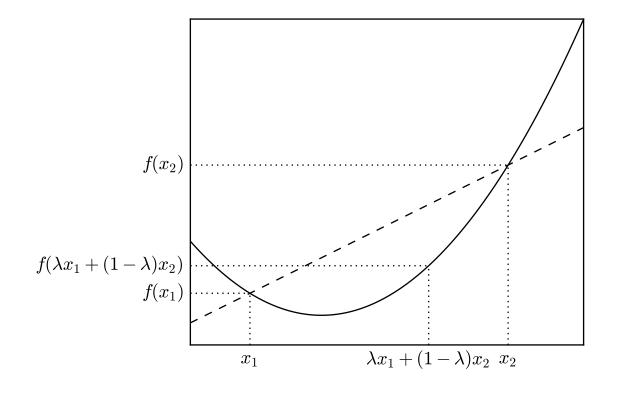
- *Quality*: Every solution is a globally optimal solution.
- *Stability*: Changing seed solutions or optimizer settings won't change the quality of the end solution.

So-called *convex* optimization problems only admit globally optimal solutions. Convex problems can be solved efficiently using deterministic methods.

Convex sets



Convex functions



$$f(\lambda x_1 + (1 - \lambda)x_2) \le \lambda f(x_1) + (1 - \lambda)f(x_2)$$

Convex optimization

 $\min f(x) \quad \text{subject to} \\ g(x) \le 0 \\ h(x) = 0$

Problem is called convex when:

- *f* is a convex function
- *g* is a convex function
- *h* is an affine function:
 - $h(x) = 0 \Leftrightarrow h(x) \le 0 \text{ and } -h(x) \le 0$.
 - *h* must be both convex and concave, i.e., affine: h(x) = ax + b.

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This is quite restrictive

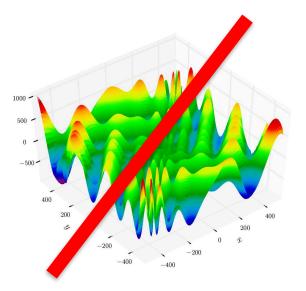
Convex problems only admit global optima.

Hydraulic modelling

- Highly nonlinear friction term in diffusive wave equation:

$$\frac{\partial H}{\partial x} + \frac{C}{R}Q^2 = 0$$

- When using many diffusive wave branches, large numbers of local minima are created. What to do?
- Linearization results in large errors; piecewise linearization results in large numbers of integer variables.





Idea: Interpolate between linearized and non-linear model.

 $M = (1 - \theta)M_L + \theta M_{NL}$

With $\theta \in [0, 1]$.

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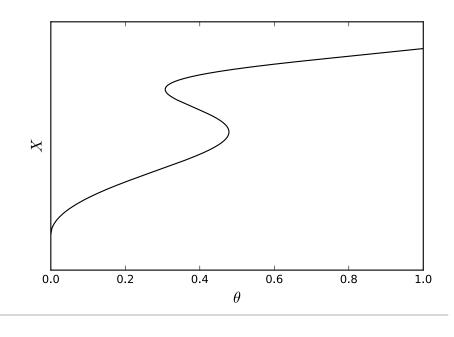




Continuation method

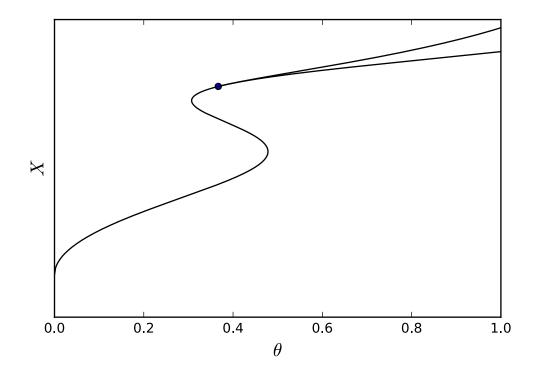
Start solving at $\theta = 0$ continuing step-by-step until $\theta = 1$.

At $\theta = 0$, the linearized problem only has global optima. We find a global optimum, and trace this through to a nonlinear solution at $\theta = 1$: Solution to nonlinear problem can be traced back to globally optimal solution of linearized problem.



Continuation method

Ongoing research challenge: Bifurcations



Suppose we have the following goals:

- Keep water levels within bounds as much as possible
- Maintain minimum spill flows for fish migration, if possible
- Apply best effort to track the generation request

Let $\{f_i : i \in I\}$ denote the set of functions encoding these goals. We have:

min $f_i \forall i \in I$ subject to $g(x) \le 0$ h(x) = 0

How to solve this?



Pareto optimality

A solution x^* of the problem

 $\min f_i \ \forall i \in I \text{ subject to} \\ g(x) \leq 0 \\ h(x) = 0$

Is *Pareto-optimal* if there is no x^{**} such that for a *j*

 $f_j(x^{**}) < f_j(x^{**})$

and for all $i \neq j$

In words: Pareto optimality implies that no goal can be improved without making another one worse.

 $f_i(x^{**}) \le f_i(x^{**})$

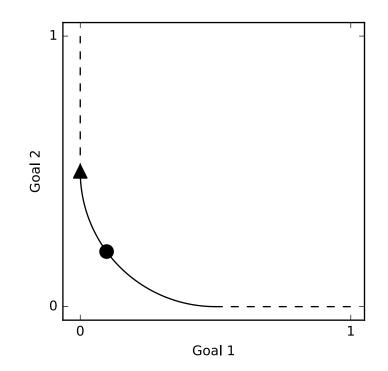
 $\min f_i \forall d$







The *Pareto front* is the set of all Pareto-optimal solutions.

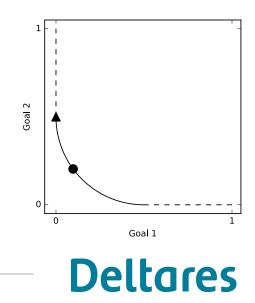


The weighting method transforms the multi-objective problem to the scalar problem

 $\min \sum_{i} \frac{\lambda_{i} f_{i}}{f_{i}} \text{ subject to}$ $g(x) \leq 0$ h(x) = 0

- Problem: How to pick the weighting factors λ_i .
- And if the weighting factors are arbitrary to a degree, then so is the solution!

Solution on Pareto front shown with a circle.



Lexicographic goal programming

In lexicographic goal programming, we transform the multi-objective problem to a <u>sequence</u> of scalar optimization problems.

First, we order our goals. For example:

- 1. Keep water levels within bounds as much as possible
- 2. Maintain minimum spill flows for fish migration, if possible

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3. Apply best effort to track the generation request

Lexicographic goal programming

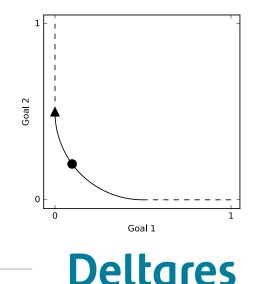
The *idea* of the algorithm is:

- 1. Minimize f_1 to yield a minimum objective value of ε_1 .
- 2. Minimize f_2 to yield ε_2 subject to the additional constraints
 - $f_1(x) = \varepsilon_1$
- 3. Minimize f_3 subject to the additional constraints
 - $f_1(x) = \varepsilon_1$

•
$$f_2(x) = \varepsilon_2$$

4. ...

Solution on Pareto front shown with an arrow.



Mixed integer: Minimum on-time pump

$$0 \le Q_t \le Q_{max}\delta_t$$

$$\delta_t - \delta_{t-1} \le 2(1 - x_t) - 1$$

$$\delta_t - \delta_{t-1} \ge -x_t$$

$$\delta_{t-1} - \delta_t \le 2(1 - y_t) - 1$$

$$\delta_{t-1} - \delta_t \ge -y_t$$

$$\sum_{i \in \{t, \dots, t+N-1\}} (x_i + y_i) \le 1$$

$$\delta_t, x_t, y_t \in \{0, 1\}$$

	δ_{t-1}	δ_t	$\delta_t - \delta_{t-1}$	x _i	<i>Y</i> _i
Pump stays off	0	0	0	0	0
Pump switched on	0	1	1	1	0
Pump switched off	1	0	-1	0	1
Pump stays on	1	1	0	0	0

Idea: Encoding logical tables using linear constraints.